## COMBINATÓRIA E TEORIA DE CÓDIGOS Exercise List 2

## 2/3/2011

Exercises 2.2 - 2.4, 2.6 - 2.11 (R. Hill)

Problem 1. a) Exercises 2.17 and 2.19 in R. Hill;

b) (Generalization of 2.19) The important familly of Reed-Muller binary codes can be obtained as follows:

 $\begin{array}{l} \forall \, r, m \in \mathbb{N}_0: \left\{ \begin{array}{l} \mathsf{RM}(0,m) = \{\vec{0},\vec{1}\} & \text{the binary repetition code with length } 2^m \\ \mathsf{RM}(m,m) = \left(\mathbb{F}_2\right)^{2^m} \\ \mathsf{RM}(r,m) = \mathsf{RM}(r,m-1) \ast \mathsf{RM}(r-1,m-1) \ , \quad 0 < r < m \end{array} \right. \end{array}$ 

where  $C_1 * C_2$  denotes the Plotkin Construction obtained from the codes  $C_1$  and  $C_2$ .

Study this family of codes by showing that the parameters of RM(r,m) are:  $n = 2^m$ ,  $M = 2^{\delta(r,m)}$ , where  $\delta(r,m) = \sum_{i=0}^r \binom{m}{i}$ ,  $d = 2^{m-r}$ .

Problem 2. a) Exercises 2.20 - 2.22 in R. Hill;

b) (Generalization of the Plotkin Bound) For q-ary codes, show that

$$A_{q}(n,d) \leq \frac{d}{d-\theta n}$$
,

where  $d > \theta n$  and  $\theta = \frac{q-1}{q}$ .